

Overview

Signed distance is commonly employed to numerically represent material interfaces with complex boundaries in multi-material numerical simulations. However, the performance of computing the signed distance field is hindered by the complexity and size of the input. Recent trends in High-Performance Computing architecture consist of multi-core CPUs and accelerators that collectively expose tens to thousands of cores to the application. Harnessing this massive parallelism for computing the signed distance field presents significant challenges. Chief among them is the design and implementation of a performance portable solution that can work across architectures. Addressing these challenges to accelerate signed distance queries is the primary contribution of this work. Specifically, in this work we employ the RAJA programming model, which provides a loop-level abstraction that decouples the loop-body from the parallel execution and insulates application developers from non-portable compiler and platform-specific directives.

RAJA Parallelization and Portability

We used RAJA's loop abstraction model to minimize the changes necessary to run our loops using different execution policies.

```

Sequential Execution
RAJA::forall<RAJA::seq_exec>(@, nnodes, signed_distance);

OpenMP Execution
RAJA::forall<RAJA::omp_parallel_for_exec>(@, nnodes, signed_distance);

Balanced OpenMP Execution using IndexSet
RAJA::IndexSet idxSet;
for(int y = 0; y <= params.ny; ++y) {
  for(int z = 0; z <= params.nz; z++) {
    int start = (params.nx + 1) * (y + z + (params.ny + 1));
    idxSet.push_back(RAJA::RangeSegment(start, start + params.nx + 1));
  }
}
typedef RAJA::IndexSet::ExecPolicy<RAJA::omp_parallel_for_segit, RAJA::simd_exec>
RAJA::forall<exec_pol> (idxSet, signed_distance);

Loop Kernel:
auto signed_distance = [=] (int index) {
  quest::Point<double, 3> pt;
  umesh->getMeshNode(index, pt.data());
  phi_store[index] =
    sd->computeDistance(pt);
  double dist = phi_store[index];
  phi_union[index] =
    std::min(dist, phi_wing[index]);
}
  
```

Conclusion

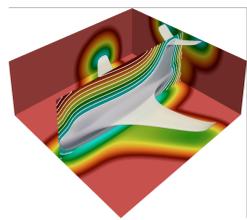
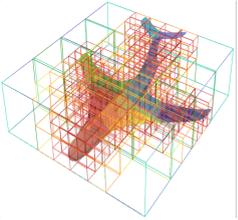
- We described a performance portable parallelization strategy for threading signed distance queries
- Our preliminary performance evaluation indicated significant load imbalances among threads
- With the RAJA IndexSet abstraction, we address the load imbalances and improve the overall performance of our algorithm

Future Work

- Use the RAJA Performance Portability Layer to
 - Evaluate our implementation on different architectures, such as machines equipped with GPUs or many-core processors, such as the Xeon Phi
 - Exploring and comparing our implementation with other programming models
- Improve our BVH queries through
 - Implementing a Surface Area Heuristic (SAH) to the BVH to achieve a more optimal BVH decomposition
 - Compaction of internal BVH data layout to optimize cache utilization and overall memory footprint

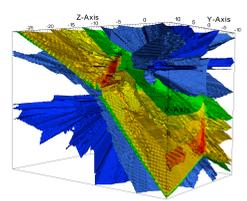
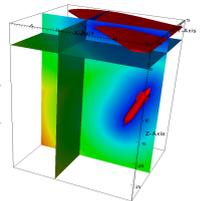
Accelerated Signed Distance

- We employed a Bounding Volume Hierarchy acceleration structure to partition the surface elements of our test configuration



BVH subdivision of a generic jet Signed Distance field on 32³ grid

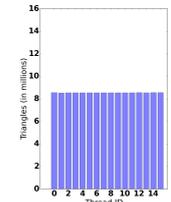
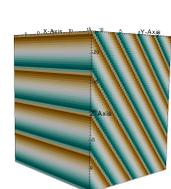
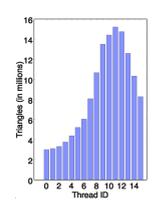
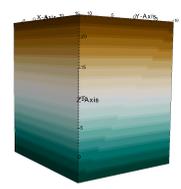
- We query the BVH subdivision to reconstruct exact signed distances from an arbitrary set of points to the surface.
- Queries to the BVH subdivision are independent and therefore ripe for the parallelization enabled by next generation architectures.



Union Signed Distance of our wing-store configuration Spatially varying cost of querying the BVH representation

Load Balancing

- We used the number of triangles tested per thread as a hardware independent workload metric
- We discovered a **510%** difference between the fastest and slowest threads in the default RAJA parallelization.
- This was reduced to **0.3%** after we applied RAJA IndexSets to change the thread partitioning.

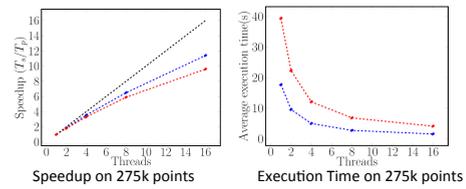


Before load balancing: Default thread mapping for a grid of 275k query points (left, colored by Thread ID) and the number of triangles (in millions) tested by each of 16 threads (right)
After load balancing: IndexSet remapping of threads to query points (left, colored by Thread ID) and the number of triangles (in millions) tested by each of 16 threads (right) on same grid

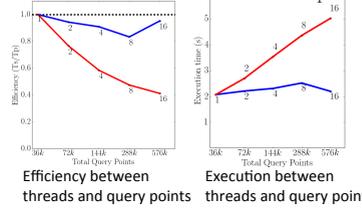
Performance Results

- After Load Balancing (blue dashed line)
- Before Load Balancing (red dotted line)
- Ideal (black dashed line)
- Strong scaling yields a **9.7x** speedup with 16 threads and **12x** after load balancing
- Load balancing increased the speedup by 23.7% and had nearly linear weak scaling efficiency

Strong Scaling



Weak Scaling



Videos and Other Media

